

Black Holes are Alpha-Bit Sup

arXiv:1706.09434, arXiv:1805.?????

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Preview

- ❑ Qubits are composite resources
- ❑ Another resource (that you have never heard of) is more fundamental than a qubit
- ❑ Quantum error correction in AdS/CFT is only approximate and bulk operators are state-dependent
- ❑ Sending qubits at the coherent information rate does not exhaust the ability of a channel to send quantum information
- ❑ There is no need to use classical bits to do entanglement-distillation, state-merging, remote state preparation, channel simulation or teleportation

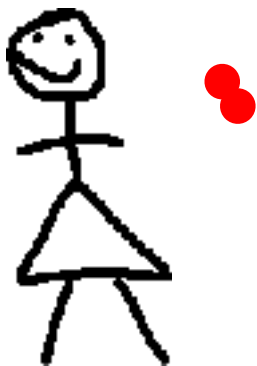
Quantum Communication Resource Inequalities

1 qubit > 1 ebit

zero-bits

1 ebit + 2 ~~ebits~~ > 1 qubit

weakened
version of qubits



Quantum Communication Resource Inequalities

$$1 \text{ cbit} < 1 \text{ qubit} > 1 \text{ ebit}$$

$$\underbrace{1 \text{ ebit}}_{\text{coherence}} + \underbrace{2 \text{ zero-bits}}_{\text{communication}} \stackrel{(a)}{=} 1 \text{ qubit}$$

weakened
version of qubits

$$m \text{ qubits} \geq 2m \text{ zero-bits} \quad ?$$

asymptotic

$$1 \text{ cbit} \stackrel{(a)}{>} 1 \text{ zero-bit} \quad ?$$

What are zero-bits?

1 zero-bit = 1 α -bit with $\alpha = 0$

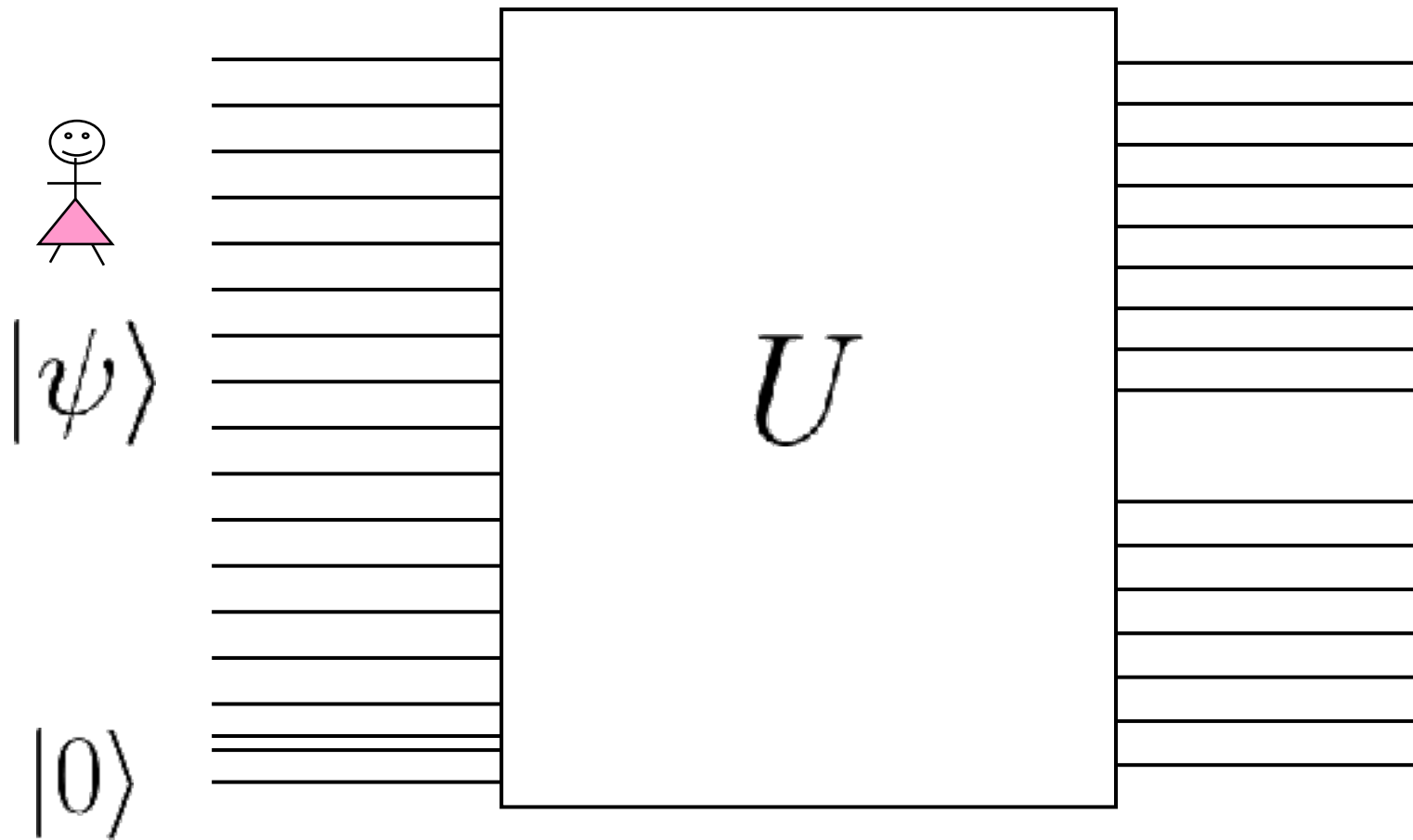
What are zero-bits?

$$d_B \gg d_E$$

B encodes the zero-bits of $|\psi\rangle$



isometric



$$\begin{aligned}
 & \forall |\psi_1\rangle, |\psi_2\rangle, \\
 & \left\| \psi_1^B - \psi_2^B \right\|_1 \approx \left\| \psi_1 - \psi_2 \right\|_1 \\
 & \forall |\psi\rangle, \\
 & \left\| \psi^E \right\|_1 \approx \omega^E
 \end{aligned}$$

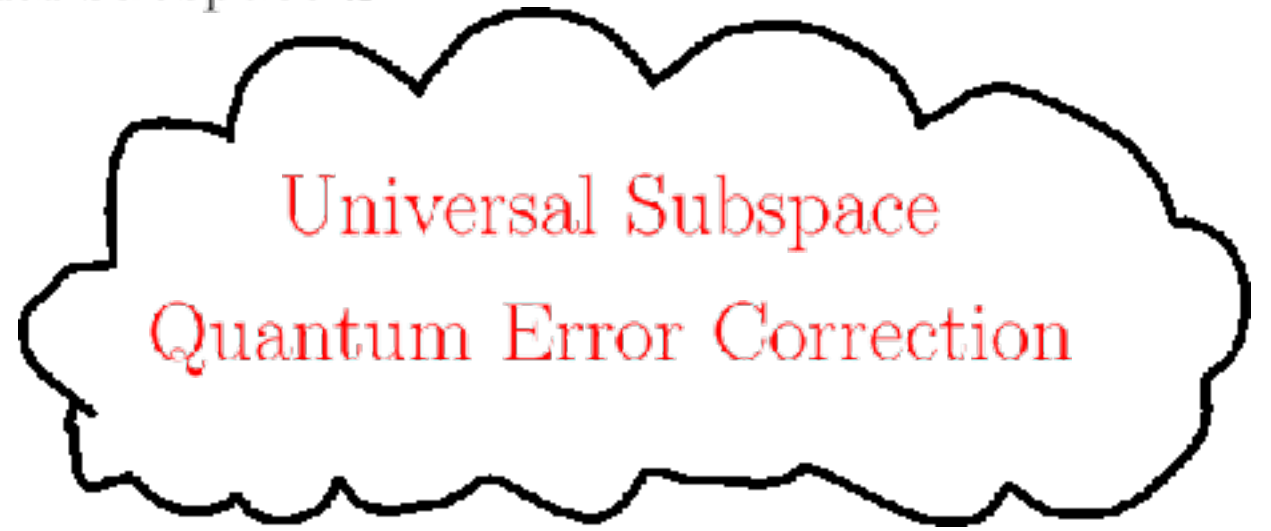
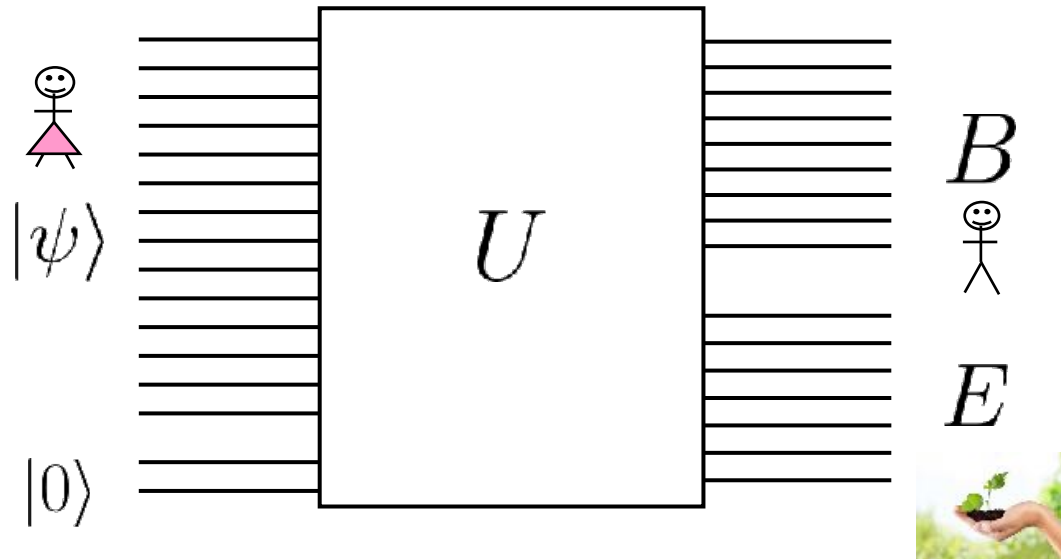


What can you do with zero-bits?

Not possible to recover the state $|\psi\rangle$ from B with no information about the state
- Error correction is not possible

However if we know $|\psi\rangle \in \text{span}(|\psi_1\rangle, |\psi_2\rangle)$ then we can determine $|\psi\rangle$

Able to error-correct any two-dimensional subspace S



Definition of zero-bits

Definition of qubits

“ n qubits”

$$\mathcal{N} : \mathcal{S}(A) \rightarrow \mathcal{S}(B)$$

$$d_A = 2^n$$

$\exists \mathcal{D}$

$$\mathcal{N} = \text{Id} \quad ?$$

What do we need to be true about the channel?

Bob can always error correct so long as error correction is possible

Definition of zero-bits

“ n zero-bits”

$$\mathcal{N} : \mathcal{S}(A) \rightarrow \mathcal{S}(B)$$

OK now what about zero-bits?

$$d_A = 2^n$$

$$\forall S \subseteq A \quad d_S = 2$$

$$\exists \mathcal{D}_S$$

$$\mathcal{D}_S \circ \mathcal{N}|_S = \text{Id}_S$$

Now Bob only has to be able to error correct any *two-dimensional subspace*

Definition of zero-bits

“ n qubits”

$$\mathcal{N} : \mathcal{S}(A) \rightarrow \mathcal{S}(B)$$

$$d_A = 2^n$$

$$\forall S \subseteq A \quad d_S = 2$$

$$\exists \mathcal{D}_S$$

$$\mathcal{D}_S \circ \mathcal{N}|_S = \text{Id}_S$$



Definition of zero-bits



“ n zero-bits”

$$\mathcal{N} : \mathcal{S}(A) \rightarrow \mathcal{S}(B)$$

$$\mathcal{N}^c : \mathcal{S}(A) \rightarrow \mathcal{S}(E)$$

[Hayden, Winter 2012]

$$d_A = 2^n$$



$$\forall |\psi\rangle \in A,$$

$$\|\mathcal{N}^c(\psi - \omega)\|_1 \leq \varepsilon$$

$$\forall S \subseteq A \quad d_S = 2$$

$$\exists \mathcal{D}_S$$

$$\|\mathcal{D}_S \circ \mathcal{N}|_S - \text{Id}_S\|_{\diamond} \leq \delta$$

Need to make definition approximate if zero-bits are to be different from qubits

$$\frac{1}{16} \delta^2 \leq \varepsilon \leq 8\sqrt{\delta}$$

Definition of alpha-bits

“ n α -bits”

“Subspace decoupling duality”

$$\mathcal{N} : \mathcal{S}(A) \rightarrow \mathcal{S}(B)$$

$$\mathcal{N}^c : \mathcal{S}(A) \rightarrow \mathcal{S}(E)$$

$$d_A = 2^n$$



$$\forall |\psi\rangle \in AR, \quad d_R = 2^{\alpha n}$$

$$\forall S \subseteq A \quad d_S \leq 2^{\alpha n} + 1 \quad \|\mathcal{N}^c \otimes \text{Id}_R (\psi - \omega \otimes \psi^R)\|_1 \leq \varepsilon$$

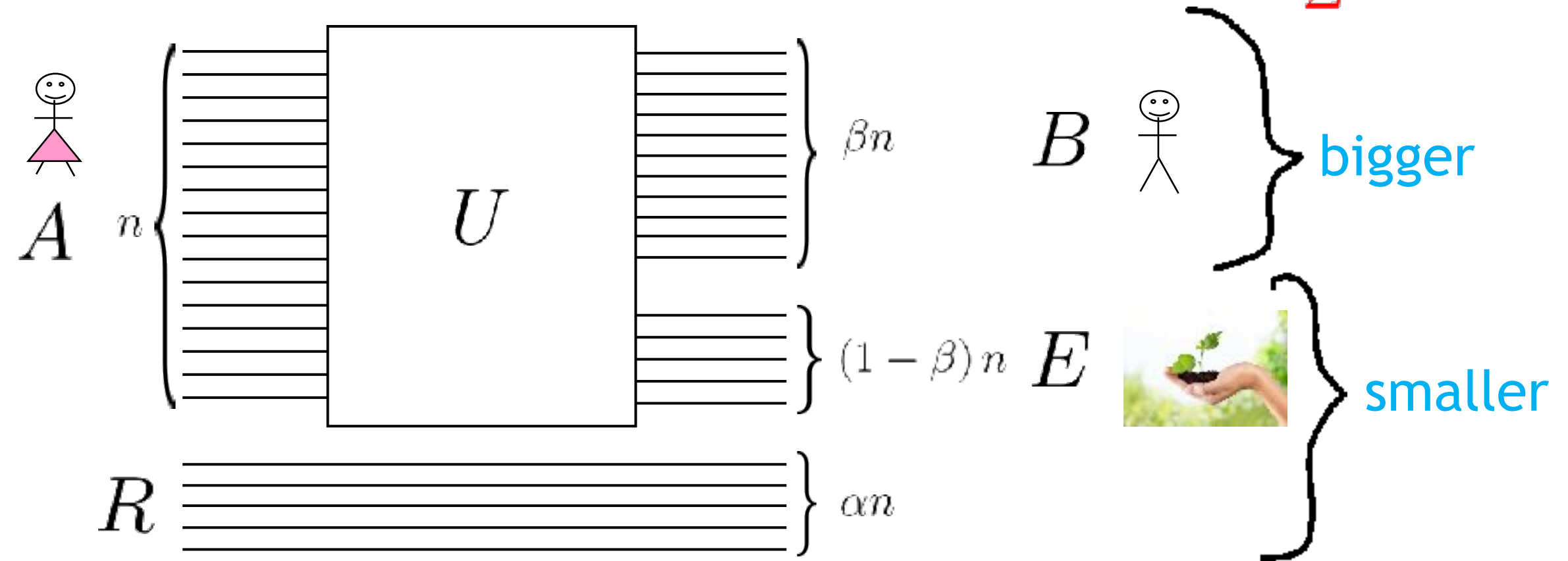
$$\exists \mathcal{D}_S \quad \alpha = 1 \Rightarrow \text{qubits}$$

$$\|\mathcal{D}_S \circ \mathcal{N}|_S - \text{Id}_S\|_{\diamond} \leq \delta$$

$$\frac{1}{16} \delta^2 \leq \varepsilon \leq 8\sqrt{\delta}$$

Transmitting alpha-bits

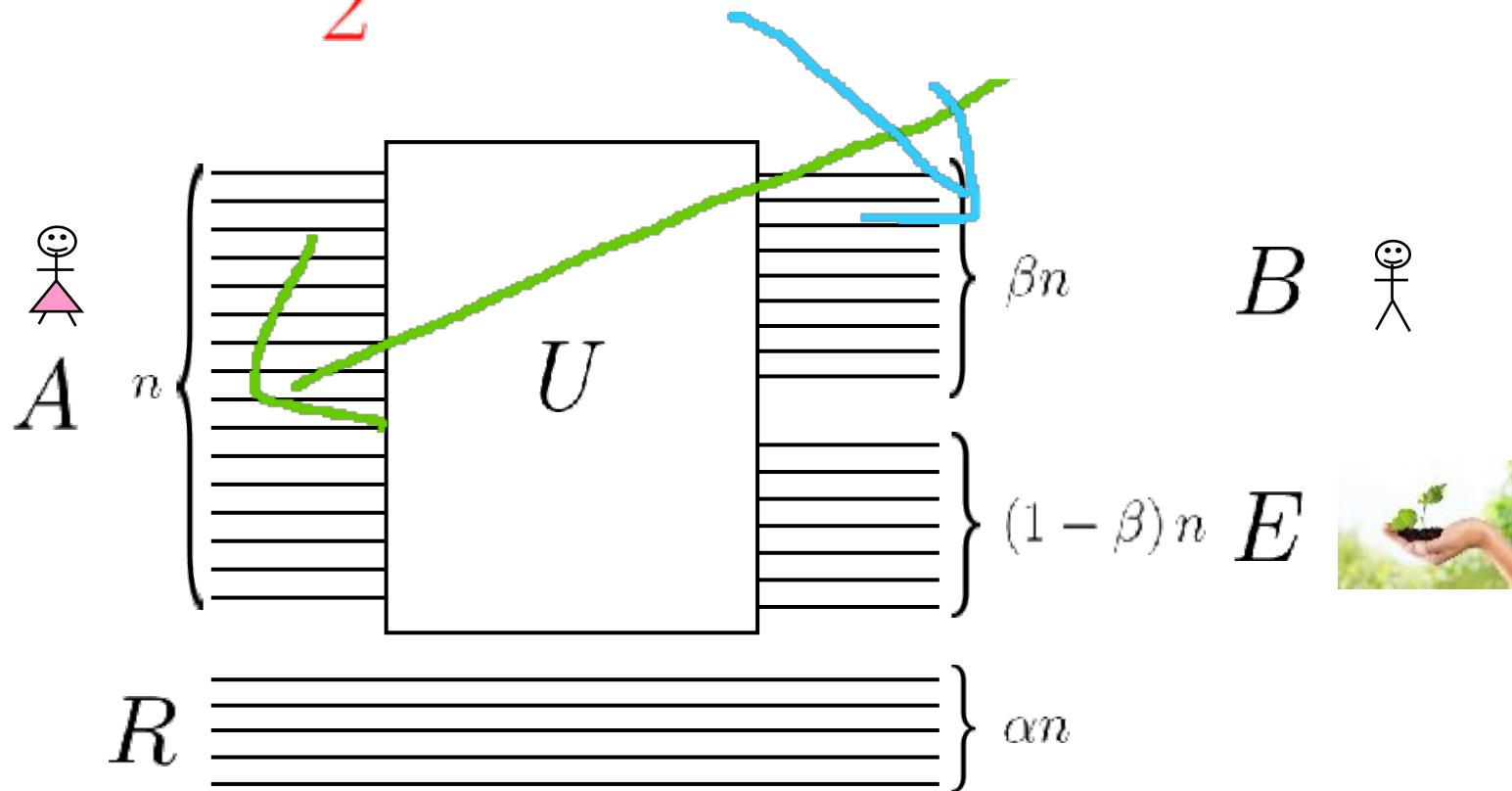
$$\beta > \frac{1 + \alpha}{2}$$



Necessary condition to send alpha-bits. Also sufficient (with some subtleties about needing to use shared randomness and block

Transmitting alpha-bits

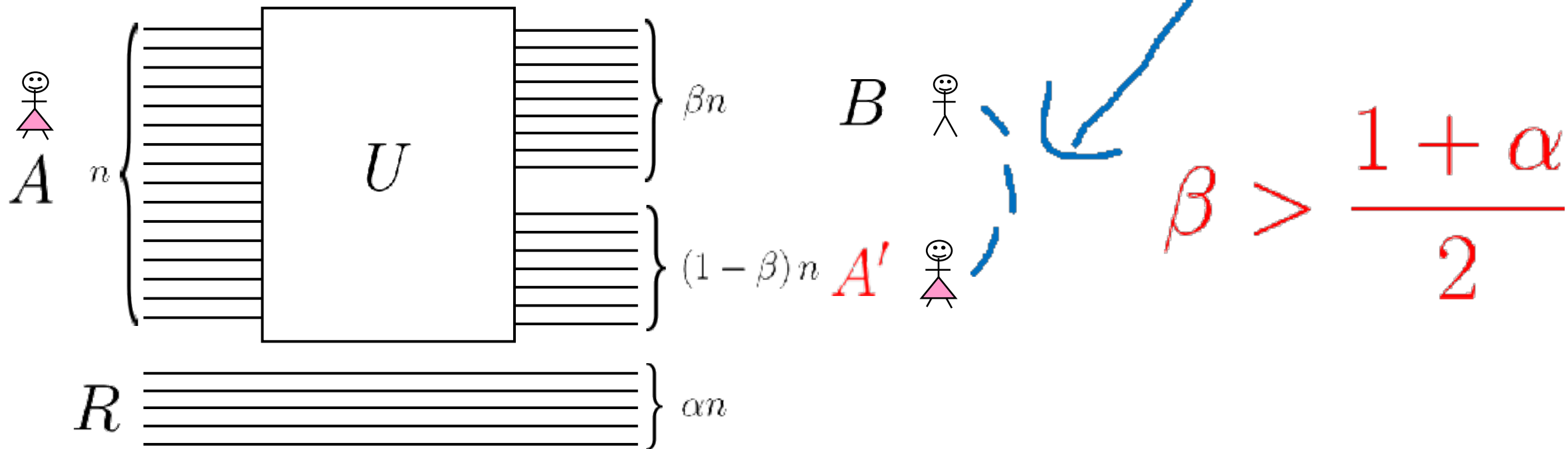
$$\frac{1 + \alpha}{2} n \text{ qubits}$$



$$\beta > \frac{1 + \alpha}{2}$$

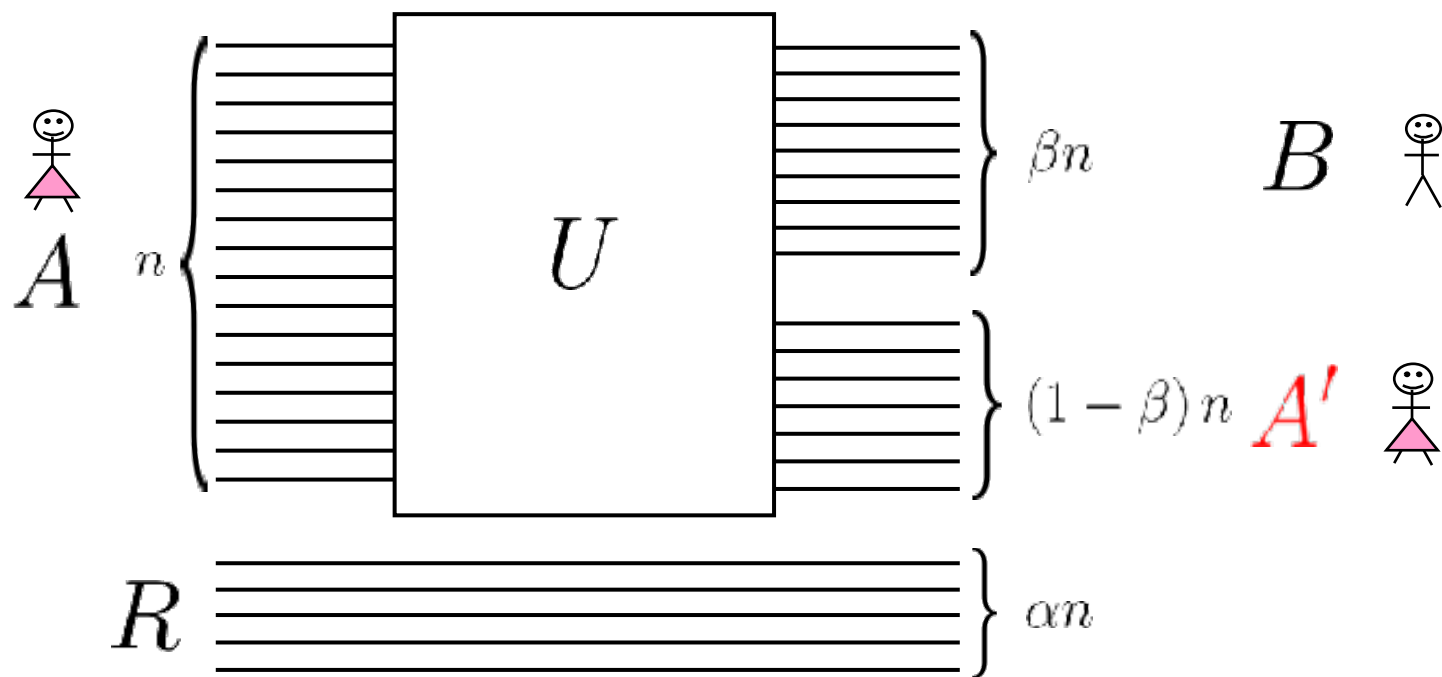
Transmitting alpha-bits

$$\frac{1 + \alpha}{2} n \text{ qubits} \geq n \alpha\text{-bits} + \frac{1 - \alpha}{2} n \text{ ebits}$$



Transmitting alpha-bits

$$(1 + \alpha) \text{ qubits} \stackrel{(a)}{\geq} 2 \alpha\text{-bits} + (1 - \alpha) \text{ ebits}$$



$$\beta > \frac{1 + \alpha}{2}$$

(Coherent) super-dense coding

1. Alice and Bob share n ebits



$$\sum_{k=0}^{2^n-1} |k\rangle_A |k\rangle_B$$



(Coherent) super-dense coding

1. Alice and Bob share n ebits
2. Alice applies an operation to her qubits
3. Alice sends her qubits to Bob

$$\sum_{k=0}^{2^n-1} e^{\frac{2\pi k r i}{2^n}} |k \oplus s\rangle_A |k\rangle_B$$

$$0 \leq r < 2^n$$

$$0 \leq s < 2^n$$



(Coherent) alpha-bit super-dense coding

1. Alice and Bob share n ebits
2. Alice applies an operation to her qubits
3. Alice sends her qubits to Bob **as α -bits**

$$\sum_{k=0}^{2^n-1} e^{\frac{2\pi k r i}{2^n}} |k \oplus s\rangle_A |k\rangle_B$$

$$0 \leq r < 2^n$$

$$0 \leq s < 2^{\alpha n}$$



Done ✓



Zero-bits and ebits as fundamental resources

All noiseless quantum resources (qubits, α -bits, cobits...) can be rewritten in terms of zero-bits and ebits

$$\text{e.g. } 1 \text{ } \alpha\text{-bit} \stackrel{(a)}{=} (1 + \alpha) \text{ zero-bits} + \alpha \text{ ebits}$$

When rewritten in this basis, the quantum resource ordering becomes the product ordering:

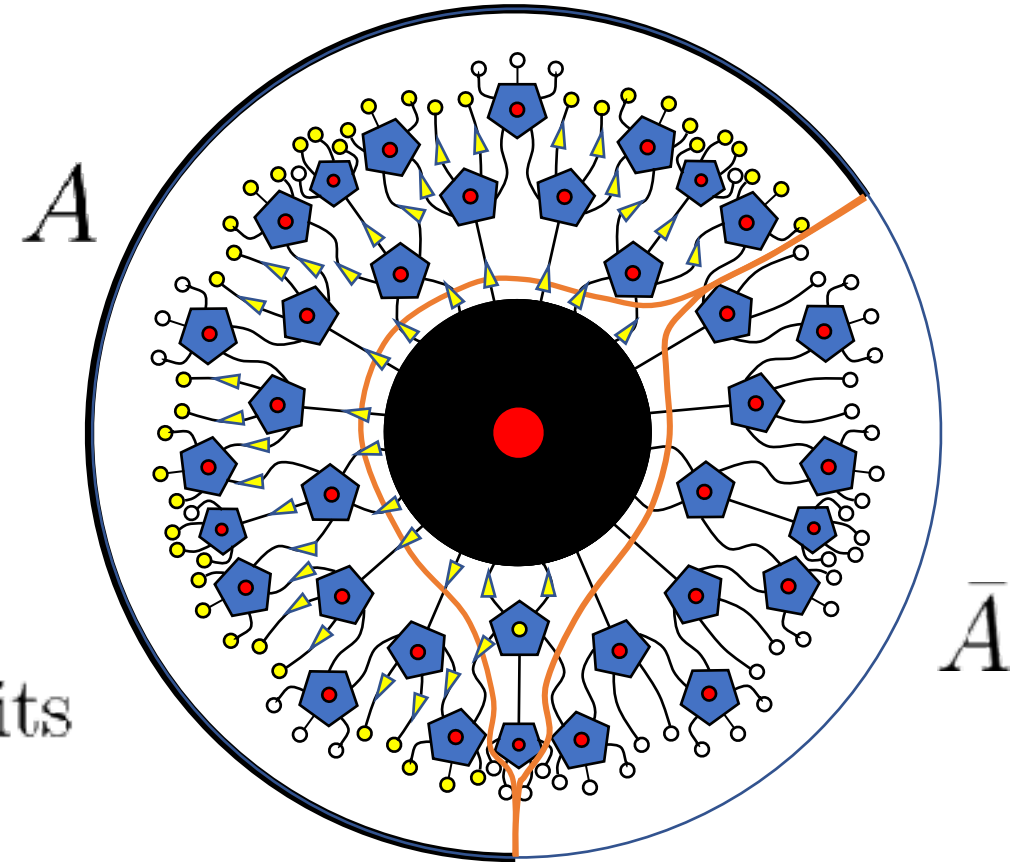
$$(a, b) \geq (a', b') \iff (a \geq a') \wedge (b \geq b')$$

Alpha-bits and Black Holes

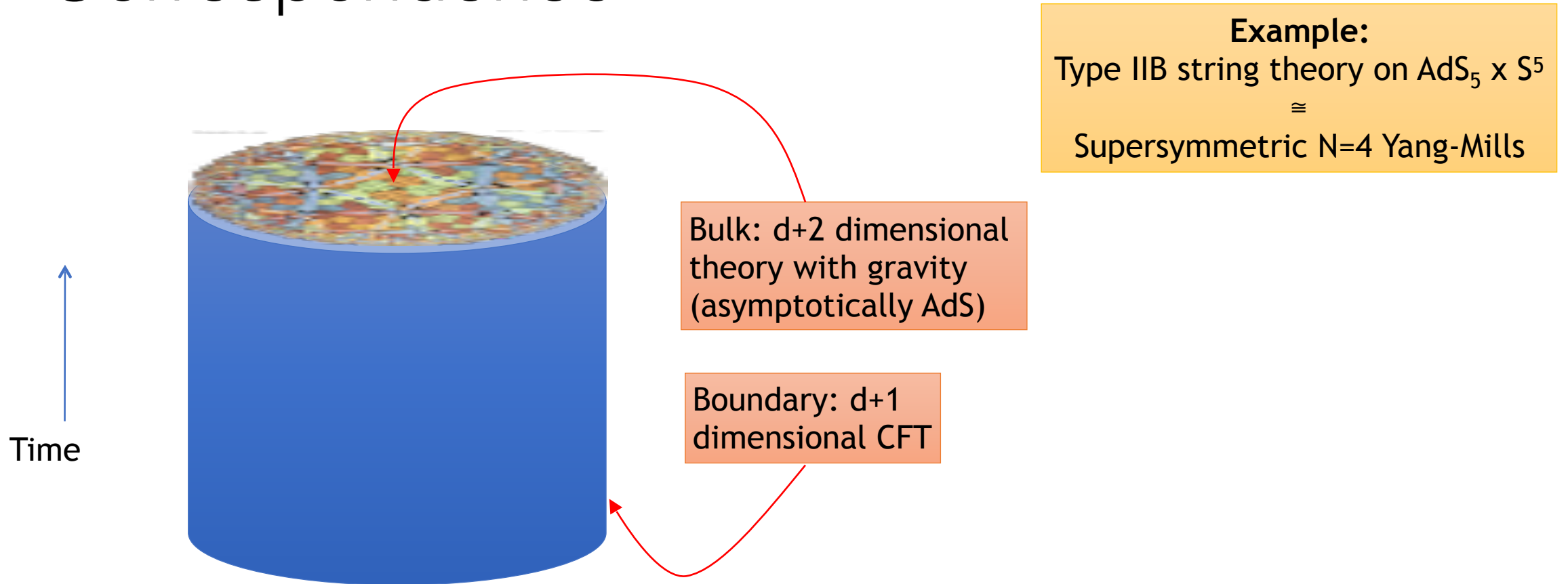
Alpha-bits arise naturally when studying black holes in AdS/CFT

Boundary subregion may encode α -bits of a bulk region

Implications: Error-correction is only approximate, reconstructed operators are state-dependent



Anti-de Sitter/Conformal Field Theory Correspondence

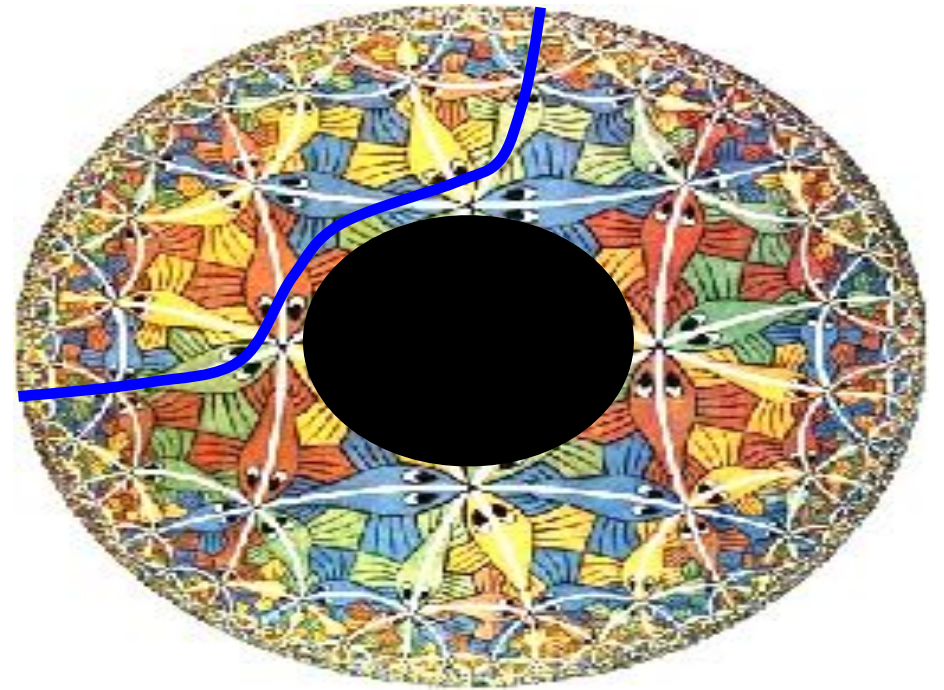


Conjecture: Equivalence of string (gravity) theory in bulk with CFT on boundary [Maldacena'97]

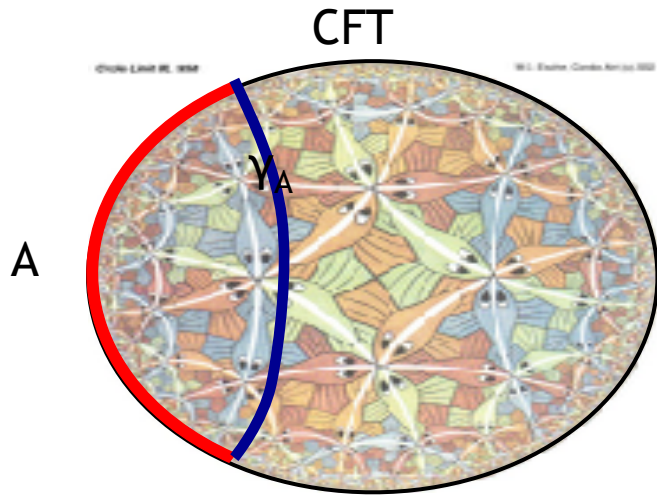
Question: How are bulk degrees of freedom encoded in the boundary?

Spatial slice of anti-de Sitter space

- Hyperbolic space
- Fish-counting metric
- Geodesics (straight lines) follow fish
- Negatively curved: sum of angles in triangle $< 180^\circ$
- Can place matter deep inside AdS



Entropy in AdS/CFT

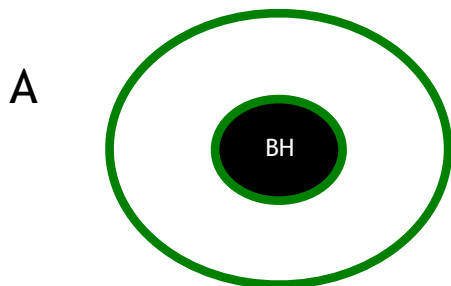


Ryu-Takayanagi proposal for bulk formula:

$$\text{QM} \curvearrowright S(A) = \frac{1}{4G_N} \min_{\gamma_A} (\text{area}(\gamma_A)) \curvearrowleft \text{GR}$$

Minimize over spatial bulk surfaces γ_A homologous to A.

Generalizes black hole entropy to wide class of spatial regions!



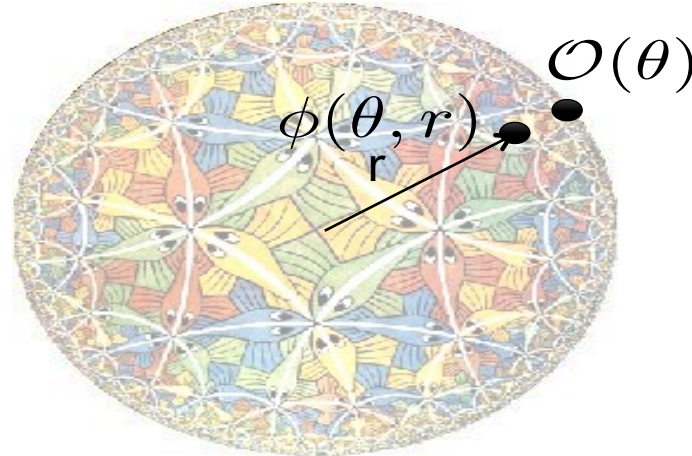
Deep message:
Entanglement structure of boundary reflects (encodes?) geometry of bulk.

- Analytical agreement in $\text{AdS}_3/\text{CFT}_2$ [RT'06]
- Satisfies strong subadditivity [Headrick-T'07]
- Proof for spherical A [Casini-Huerta-Myers'11]
- General explanation [Lewkowycz-Maldacena'13]

Relating bulk and boundary observables

Boundary in terms of bulk: *extrapolate*

$$\mathcal{O}(\theta) = \lim_{r \rightarrow \infty} r^\Delta \phi(\theta, r)$$



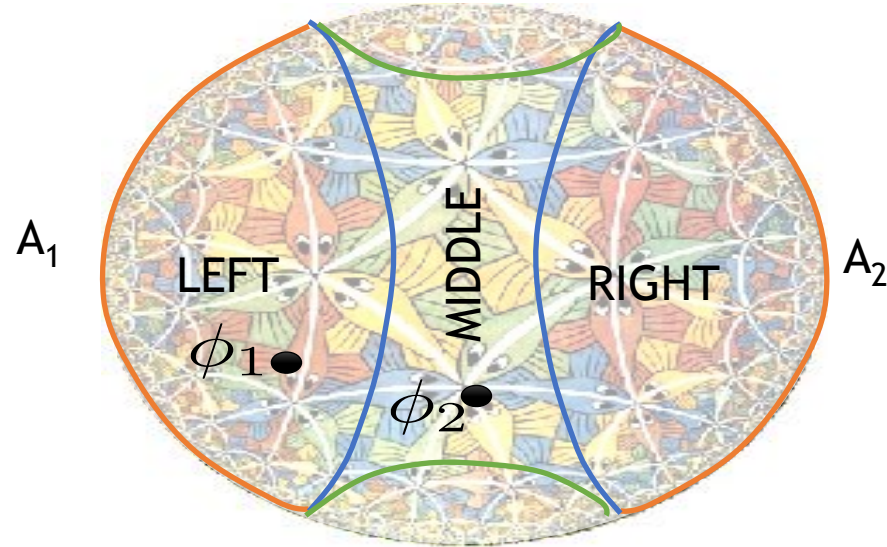
Bulk in terms of boundary: *smearing*

$$\phi(\theta, r) = \int K(\theta, r; \tilde{\theta}, \tilde{t}) \mathcal{O}(\tilde{\theta}, \tilde{t}) d\tilde{\theta} d\tilde{t}$$

K arises from solving some PDE's (Green's fn for classical bulk field equations)

Don't always need the whole boundary to reconstruct a given $\phi(\theta, r)$.

The causal and entanglement wedges



$$S(A) = \frac{1}{4G_N} \min_{\gamma_A} (\text{area}(\gamma_A))$$

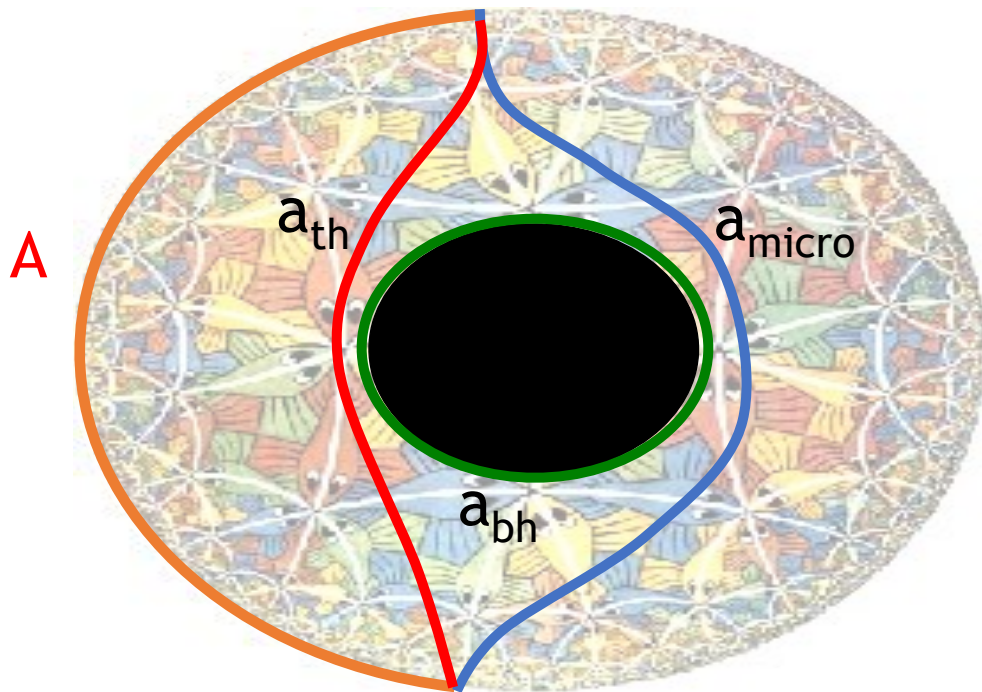
Causal wedge: LEFT + RIGHT

Entanglement wedge: LEFT + MIDDLE + RIGHT

Hamilton, Kabat, Lifschytz, Lowe:
Can reconstruct all bulk operators in the causal wedge

Quantum information arguments:
Can reconstruct all bulk operators in the full entanglement wedge

Decoding black hole microstates



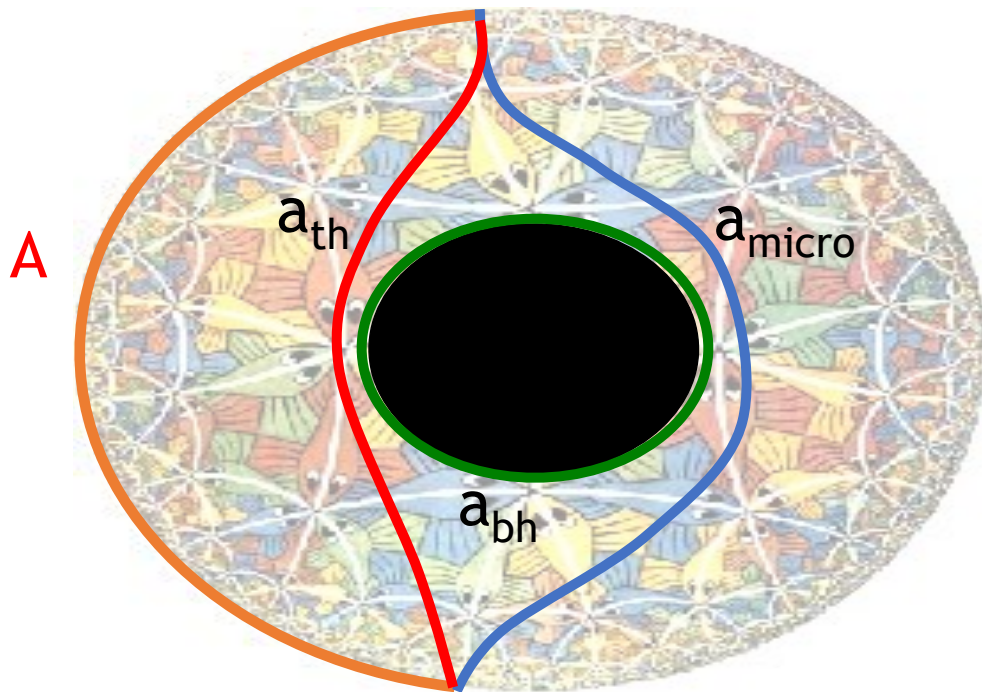
Pure state: a_{micro} is minimal
BH observables are mapped to A

Thermal state:
Minimizing with homology constraint
 $a_{micro} + a_{bh} > a_{th}$

BH is outside A's thermal entanglement wedge:
BH observables not mapped to A

Location of entanglement wedge depends
on how uncertain decoder is about the microstate

Decoding black hole microstates



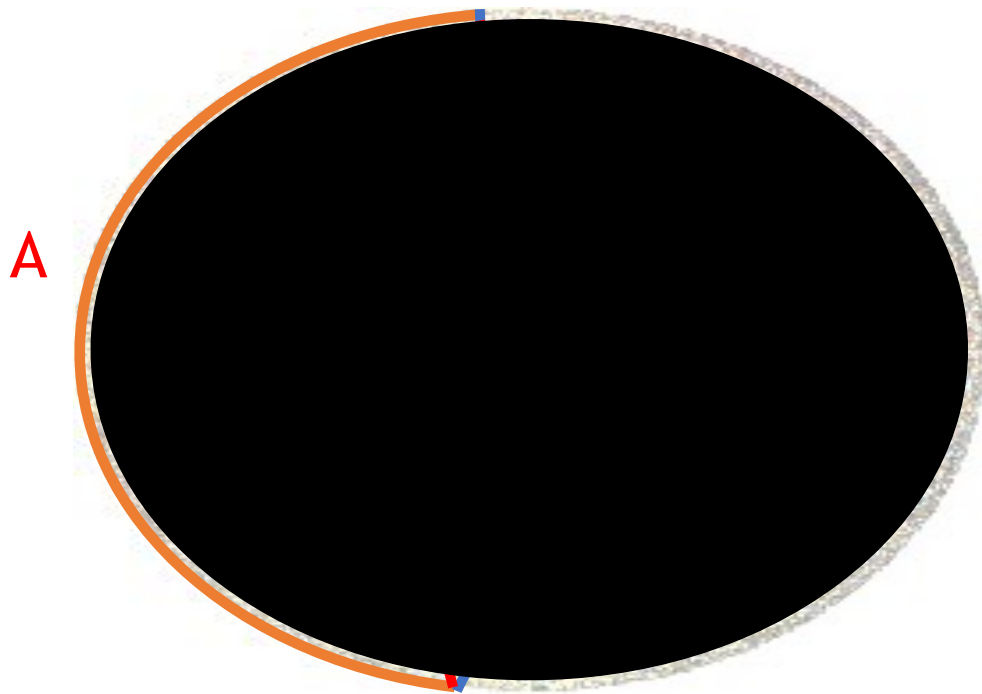
Consider mixture of BH microstates
Entropy: $r = \alpha a_{bh}$

Minimizing with homology constraint
Choose α such that
$$a_{micro} + \alpha a_{bh} < a_{th}$$

a_{micro} is minimal:
BH observables mapped to A

Any mixture of microstates will do: universal subspace quantum error correction!
 A contains the black hole's α -bits!

Black holes are alpha-bit sup



Consider mixture of BH microstates
Entropy: $r = \alpha a_{bh}$

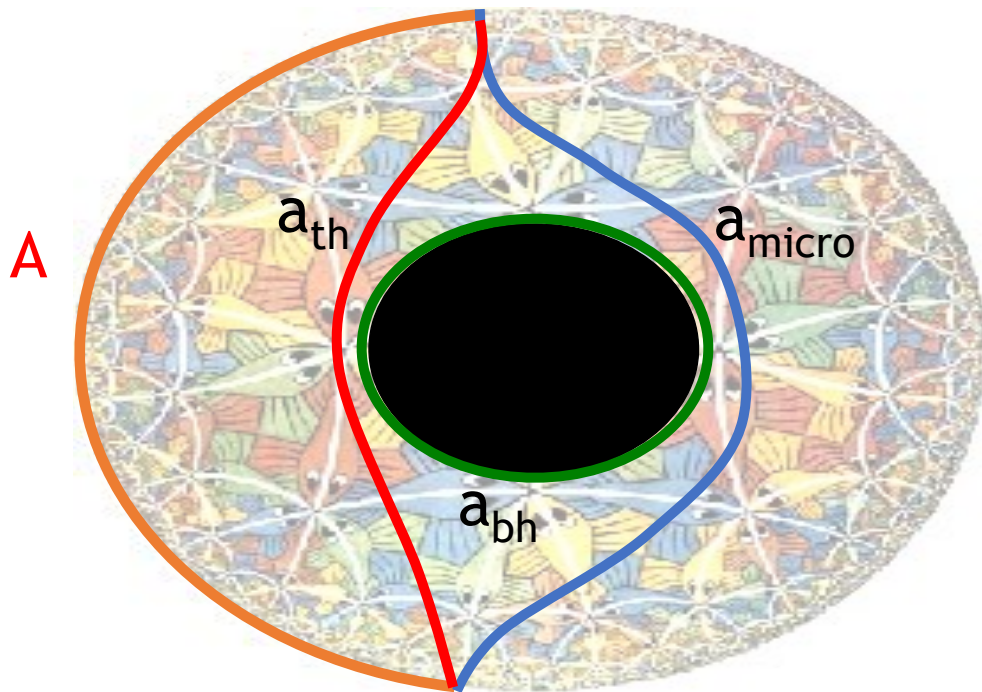
Minimizing with homology constraint
Choose α such that
 $\alpha < (a_{th} - a_{micro})/a_{bh}$

$$a_{th} = A = x a_{bh} \quad a_{micro} = A^c = (1-x) a_{bh}$$

$x > (1+\alpha)/2$: saturates α -bit capacity!

Large black holes in AdS/CFT give explicit (but #\$\$%! complicated) optimal α -bit codes for all α

Consequences



- Bulk to boundary mapping must be approximate
 - Otherwise, no difference between alpha-bits and qubits
 - Corrections $\exp(-O(S_{bh}))$
 - Nonperturbative?
- Bulk to boundary mapping is state-dependent, even in a fixed geometry

Summary

- Alpha-bits quantify asymptotically distinct forms of approximate quantum error correction
- Qubits are composite resources
 - Ebits and 0-bits are the fundamental independent resources of correlation and communication
- Alpha-bits arise naturally in AdS/CFT
 - Inequivalence of alpha-bits and qubits implies AdS \rightarrow CFT error correction is approximate
- 0-bits can substitute for cbits in entanglement distillation, teleportation, state merging, channel simulation, ...
- Amortized capacity is singular at $\alpha=1$ (quantum capacity point)

Thank you