Loop Amplitudes in QCD, Geometry and Ansätze

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The Future of Precision at the LHC



- ▶ LHC experiments will increase in statistics over next 2 decades.
- Experimental precision will reach 1% level for many observables.
- Precise theoretical predictions needed to match experimental error.

Precise Perturbative Predictions

Perturbation theory is our major tool for making predictions.

As a proton machine, the LHC is dominated by QCD.

$$\sigma[\alpha_{5}] \sim \sigma_{\mathsf{LO}} + \alpha_{5} \delta \sigma_{\mathsf{NLO}} + \alpha_{5}^{2} \delta \sigma_{\mathsf{NNLO}} + \mathcal{O}(\alpha_{5}^{3}).$$

► At LHC energies, $\alpha_S \sim 0.1$, \Rightarrow we require NNLO QCD. (+ parton-shower + resummation).

$$A_{5} = \begin{bmatrix} \mathbf{m} \\ \mathbf{m}$$

 \Rightarrow Need two-loop scattering amplitudes.

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Progress in $2 \rightarrow 3$ calculations sparked by appropriate techniques.

Introduction	Modern Techniques	Geometry of \mathcal{C}_k	Applications	Conclusions
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Outline



▶ Part 2: A Geometric Approach to Amplitude Construction.

Introduction	Modern Techniques	Geometry of \mathcal{C}_k	Applications	Conclusions
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Part 1: Modern Techniques For QCD Amplitudes

Anatomy of a Loop Amplitude

• Textbook: A loop amplitude is the sum of Feynman diagrams.

$$\mathcal{A}^{(2)}(p_1,\ldots,p_5)= {}^{\infty}_{\mathbb{R}} \mathcal{B}^{\infty}_{\mathbb{R}} \mathcal{B}^{\infty}_{\mathbb{R}} + \mathcal{O}(10000) ext{ diagrams}.$$

Each (gauge-dependent) diagram leads to Feynman integrals:

$$\sum_{p_{1}}^{D_{1}} \sum_{D_{2} \atop D_{2}}^{D_{5}} \sum_{D_{4}}^{D_{5}} \sum_{D_{6}}^{D_{6}} D_{6} = \int [d^{D}\vec{\ell}] \frac{N(\ell_{1},\ell_{2})}{D_{1}D_{2}D_{3}D_{4}D_{5}D_{6}D_{7}D_{8}}.$$

• QCD numerators N are polynomials of $\vec{\ell}$ with 1000s of terms.

Function of many variables with complex branch cut structure.

The Modern Perspective and Challenges

Aim: efficient, stable evaluation of amplitude over phase space.

• Too many integrals \Rightarrow write $A^{(2)}$ in terms of master integrals:

$$A^{(2)}(p_1,\ldots,p_5) = \sum_k \underbrace{\mathcal{C}_k(p_1,\ldots,p_5)}_{\text{rational functions}} \underbrace{\mathcal{I}_k(s_{12},s_{23},s_{34},s_{45},s_{51})}_{\text{master integrals}}.$$

• Coefficients C_k are process specific, integrals \mathcal{I}_k are universal.

Challenges

- \mathcal{I}_k : analytic complexity \Rightarrow efficient, stable evaluation difficult.
- \triangleright C_k : intermediate expression swell, poorly understood structure.

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Master Integrals

$\sum_{k} C_{k}(p_{1},\ldots,p_{5}) \mathcal{I}_{k}(s_{12},s_{23},s_{34},s_{45},s_{51})$

Differential Equations and Master Integrals

Integrals are special functions depending on many variables.

$$\mathcal{I}_{k} \sim \int [d^{D}\vec{\ell}] \frac{N(\ell_{1},\ell_{2})}{D_{1}D_{2}D_{3}D_{4}D_{5}D_{6}D_{7}D_{8}} \qquad D_{i} = \Big(\sum_{j} a_{ij}\ell_{j} + \sum_{j} b_{ij}p_{j}\Big)^{2} - m_{i}^{2}.$$

Analytic structure is manifested by differential equations.

$$\mathrm{d}\mathcal{I}_k = \mathbf{M}_{kl}(\epsilon, \vec{s})\mathcal{I}_l.$$

[Gehrmann, Remiddi '01]

► For (pure) polylogarithmic integrals, M simplifies: [Henn '13]



Powerful DE construction/initial value approaches exist.

[Abreu, BP, Zeng '18] [Abreu, Ita, Moriello, BP, Tschernow, Zeng '20]

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Numerically Solving the Differential Equation

Aim: numerically evaluate \mathcal{I}_j as ϵ series, $\mathcal{I}_j(\vec{p}, \epsilon) = \sum_{k=-4}^{0} d_{j,k} \epsilon^k h_k + \mathcal{O}(\epsilon)$.

Pentagon functions: [Gehrmann, Henn, Io Presti '18] Dedicated iterated integral code: $h_k \sim \int_0^1 d \log(W_n[t_n]) \cdots \int_0^{t_2} d \log(W_1[t_1]).$

Very efficient for five-point.

Series expansions: [Moriello '19]

- ▶ Patch together h_k from power series: h_k ~ ∑_{j1,j2} (t − t₀)^{j1/2} log(t − t₀)^{j2}.
 ▶ Implementations: DiffExp/AMFlow. [Hidding '20] [Liu, Ma '22]
 - Promising for elliptic five-point.



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Master Integrals – State of the Art (Two-Loop Five-Point)

Five-point Massless



[Papadopoulos, Tommasini, Wever '15] [Gehrmann, Henn, Io Presti '18] [Abreu, Page, Zeng '18] [Chicherin, Gehrmann, Henn, Io Presti, Mitev, Wasser '18] [Abreu, Dixon, Herrmann, Page, Zeng '18] [Chicherin, Gehrmann, Henn, Wasser, Zhang, Zoia '18] [Gehrmann, Henn, Io Presti '18] [Chicherin, Sotnikov '20]



[Abreu, Ita, Moriello, Page Tschernow, Zeng '20] [Abreu, Ita, Page, Tschernow '21] [Papadopoulos, Tommasini, Wever 15]

Five-point One-Mass

[Canko Papadopoulos, Syrrakos 20] [Kardos, Papadopoulos, Smirnov, Syrrakos, Wever 19] [Chicherin, Sotnikov, Zoia '21]

[Abreu, Chicherin, Ita, Page, Sotnikov, Tschernow, Zoia. to appear]

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Integral Coefficients

$\sum_{k} C_{k}(p_{1},\ldots,p_{5}) \mathcal{I}_{k}(s_{12},s_{23},s_{34},s_{45},s_{51})$

The Analytic Reconstruction Approach

► Analytic coefficients built from numerical samples via Ansatz.

$$\mathcal{C}_k(p_1,\ldots,p_n) = \sum_{j=1}^{N_k} c_{jk} \mathfrak{a}_{jk}(p_1,\ldots,p_n), \qquad c_{jk} \in \mathbb{Q}.$$

Numerical evaluations provide constraints on unknown c_{jk}.

$$(p_1^{(0)},\ldots,p_n^{(0)})\longrightarrow \qquad \longrightarrow \mathcal{C}_1(p_1^{(0)},\ldots,p_n^{(0)}).$$

- Made practical by finite field methods (working modulo p). [Schabinger, von Manteuffel '14; Peraro '16]
- Sidesteps complex algebra only intermediate numerics!

Ansätze and [Known] Structure of QCD Coefficients

▶ Re-express C_a as rational function of Mandelstam variables.

$$\mathcal{C}_{a}(\varepsilon_{1}, p_{1}, \dots, \varepsilon_{n}, p_{n}) = \underbrace{\phi(\varepsilon_{1}, p_{1}, \dots, \varepsilon_{n}, p_{n})}_{\text{phase weight}} \underbrace{\frac{\mathcal{N}_{a}\left[s_{ij}, \epsilon(p_{i}, p_{j}, p_{k}, p_{l})\right]}{\prod_{\alpha} \mathcal{W}_{\alpha}(s_{ij})^{q_{a\alpha}}}.$$

W_α are symbol letters, from differential equation.
 [Abreu, Dormans, Febres Cordero, Ita, BP '18]

 $\mathcal{N} \text{ is polynomial in Lorentz invariants, e.g. (5-point massless): }$ $\mathcal{N} = \underbrace{\mathcal{N}_+(s_{12}, s_{23}, s_{34}, s_{45}, s_{51})}_{\text{polynomial}} + \epsilon(p_1, p_2, p_3, p_4) \underbrace{\mathcal{N}_-(s_{12}, s_{23}, s_{34}, s_{45}, s_{51})}_{\text{polynomial}}.$

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Gathering Numerical Data: Numerical Unitarity

Aim: Numerically reduce Feynman diagram sum to masters.

$$\sum_{\substack{\substack{\substack{m \in \mathcal{M}}\\m \in \mathcal{M}}}} \mathcal{M}_{m} + \mathcal{O}(10000) \text{ diagrams} = \int [\mathrm{d}^{D} \vec{\ell}] \, \mathcal{A}(\vec{\ell}) = \sum_{j} \mathcal{C}_{j} \mathcal{I}_{j}.$$

Unitarity lets us build integrand from product of trees.

$$\operatorname{Disc}\left[A\right] = \operatorname{Disc}\left[A\right] = \operatorname{Di$$

Non-uniqueness of integrand captured by surface terms σ .

$$\mathcal{A}(ec{\ell})\sim\mathcal{A}(ec{\ell})+\sigma(ec{\ell}) \qquad ext{where} \qquad \int [\mathrm{d}^Dec{\ell}]\,\sigma(ec{\ell})=0.$$

Surface terms built with algebraic geometry approaches. [Gluza et al '11; Ita '15; Larsen et al '15, Abreu, Febres Cordero, Ita, Jaquier, BP, Zeng '17] Complexity of Analytic Reconstruction

• Currently $T_{\text{reconstruct}} = \mathcal{O}(\text{laptop}), T_{\text{evaluate}} = \mathcal{O}(\text{cluster}).$

Number of evaluations for recent QCD helicity amplitudes:



Lesson: Need to better understand rational functions in QCD.

Modern Techniques	Geometry of \mathcal{C}_k	Applications	Conclusions
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Part 2: Structure of C_k from Algebraic Geometry (based on [arXiv:2203.04269] w/ Giuseppe De Laurentis)

Spinor Helicity Methods

Amplitudes conventionally represented in momentum space.

$$\sum_{p_n \ p_{n-1}}^{p_1} \sum_{p_{n-1}}^{p_2} p_3 \qquad \rightarrow \qquad \sum_{i=1}^n p_i^\mu = 0, \qquad p_i^2 = m_i^2.$$

Yang-Mills amplitudes compactly expressed in spinor variables.

$$p_i^2 = 0 \qquad \Rightarrow \qquad p_{i\mu}\sigma^{\mulpha\dot{lpha}} = \lambda_i^lpha \tilde{\lambda}_i^{\dot{lpha}}.$$

p_i is automatically on-shell, momentum conservation becomes

$$\sum_{i=1}^n \lambda_i^{\alpha} \tilde{\lambda}_i^{\dot{\alpha}} = 0.$$

Invariant quantities represented in terms of spinor brackets.

$$\langle ij \rangle = \lambda_i^{\alpha} \lambda_{j\alpha}, \qquad [ij] = \tilde{\lambda}_{i\dot{\alpha}} \tilde{\lambda}_j^{\dot{\alpha}}.$$

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Constraints from Singular Limits

• Consider six-point one-quark line amplitude $\mathcal{A}_{q^+g^+g^+\overline{q}^-g^-g^-}^{\text{tree}}$.

$$\mathcal{A} = rac{\mathcal{N}^{*}}{\langle 12
angle \langle 23
angle \langle 34
angle [45] [56] [61] s_{345}}$$

 $^*\mathcal{N}$ is a degree 6 polynomial in spinor brackets.

• Doubly singular behaviour can be used to constrain A.

[De Laurentis, Maître '19]

• Consider setting $\langle 12 \rangle$, $\langle 23 \rangle$ are small:

$$\begin{array}{ccc} \lambda_{2}^{\alpha} \sim \epsilon & \Rightarrow & \mathcal{A} \sim \epsilon^{-2} \\ \langle 12 \rangle \sim \langle 23 \rangle \sim \langle 13 \rangle \sim \epsilon & \Rightarrow & \mathcal{A} \sim \epsilon^{-1}. \end{array}$$

How do we systematically build an Ansatz with this behavior?

Challenges to Overcome

Aim Build Ansatz $\{a_1, \ldots, a_N\}$ matching singular behaviour.

$$\mathcal{C}_i(\lambda, ilde{\lambda}) = rac{\mathcal{N}_i(\lambda, ilde{\lambda})}{\prod_{j=0}^{n-1}\mathcal{D}_j(\lambda, ilde{\lambda})^{q_{ij}}} = \sum_{j=1}^N c_{ij}\mathfrak{a}_j(\lambda, ilde{\lambda}), \qquad c_{ij}\in\mathbb{Q}.$$

Strategy: Find polynomial space enjoying \mathcal{N} 's vanishing behavior. **Difficulties**:

- ► How do we handle linear dependencies in spinor space?
- ► How to systematically study (define) all singular surfaces?
- How to algorithmically turn vanishing constraints into Ansatz?

Algebraic Geometry Warmup



Algebraic Geometry and Spinors

• Workhorse: ring of polynomials in spinor components $-S_n$.

$$S_n = \mathbb{C}[\lambda_1^{lpha}, \ldots, \lambda_n^{lpha}, \tilde{\lambda}_1^{\dot{lpha}}, \ldots, \tilde{\lambda}_n^{\dot{lpha}}].$$

• "Ideal" is set of combinations of polynomials $q_i \in S_n$.

$$\langle \underbrace{q_1,\ldots,q_k}_{\text{generators}} \rangle_{S_n} = \left\{ \sum_{i=1}^k a_i q_i, \ a_i \in S_n \right\}.$$

• "Variety" associated to ideal is set of points satisfying $q_i = 0$.

$$V\Big(\langle q_1,\ldots,q_k\rangle_{S_n}\Big)=\Big\{(\lambda,\tilde{\lambda})\in\mathbb{C}^{4n}:q_i(\lambda,\tilde{\lambda})=0 ext{ for } 1\leq i\leq k\Big\}.$$

Momentum Conservation and the Quotient Ring

• Momentum conservation ideal J_{Λ_n} . Associated variety $V(J_{\Lambda_n})$.

$$J_{\Lambda_n} = \left\langle \sum_{i=1}^n \lambda_{i\alpha} \tilde{\lambda}_{i\dot{\alpha}} \right\rangle_{S_n}, \qquad V(J_{\Lambda_n}) \sim$$

▶ Polynomials differing by J_{Λ_n} element are physically equivalent.

$$p \sim q \iff p - q \in J_{\Lambda_n}$$

Simple example at four points:

 $\lambda_{1}^{\alpha}\langle 12\rangle = -\lambda_{3}^{\alpha}\langle 32\rangle - \lambda_{4}^{\alpha}\langle 42\rangle \quad \Leftrightarrow \quad \lambda_{1}^{\alpha}\langle 12\rangle + \lambda_{3}^{\alpha}\langle 32\rangle + \lambda_{4}^{\alpha}\langle 42\rangle \in J_{\Lambda_{4}}.$

▶ Polynomial spinor functions on $V(J_{\Lambda_n})$ live in quotient ring^{*}:

$$R_n = S_n/J_{\Lambda_n}.$$

*Systematically organized by Gröbner basis techniques.

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	Modern Techniques	Geometry of \mathcal{C}_k	Applications	Conclusions
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Algebra and Geometry of Singular Varieties

Introduction M	odern Techniques	Geometry of C_k	Applications	Conclusions
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Singular Varieties

• Denominator of coefficients C_i factorizes.

$$\mathcal{C}_i(\lambda, \tilde{\lambda}) = rac{\mathcal{N}_i(\lambda, \tilde{\lambda})}{\prod_{j=0}^{n-1} \mathcal{D}_j(\lambda, \tilde{\lambda})^{q_{ij}}}.$$

Example D_j , five point two-loop massless amplitudes:

$$ec{\mathcal{D}}^{(5 ext{-point})} = \{ \langle 12
angle, \langle 13
angle, \dots, [12], [13], \dots, [1|2+3|1
angle, \dots \}$$

We will consider varieties where two denominators are singular:

$$\mathcal{D}_i = \mathcal{D}_j = 0.$$

Must control branching structure and vanishing functions.

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Branching of Surfaces Defined by Polynomials

▶ When we intersect surfaces, we may have multiple branches.



Our double denominator zero surface has two branches:

$$\langle 12\rangle = \langle 23\rangle = 0 \quad \Leftrightarrow \quad \langle 12\rangle = \langle 23\rangle = \langle 13\rangle = 0 \quad \text{or} \quad \lambda_2^\alpha = 0.$$

Compute branchings with primary decomposition techniques.
 [De Laurentis, BP '22], see also [Zhang '12].

Polynomials that Vanish on a Variety

Polynomials vanishing on all points of variety U form ideal

$$I(U) = \Big\{ q \in S_n \quad ext{where} \quad q(x) = 0 \quad ext{for all } x \in U \Big\}.$$

• Consider if \mathcal{N}_i vanishes to order k on U,

$$\mathcal{N}_i(x^{(\epsilon)}) = \mathcal{O}(\epsilon^k), \quad ext{where} \quad |x-x^{(\epsilon)}| \leq \epsilon \quad ext{and} \quad x \in U.$$

Zariski-Nagata: \mathcal{N}_i belongs to k^{th} "symbolic power" $I(U)^{\langle k \rangle}$.

$$I(U)^k = I(U)^{\langle k \rangle} \cap J_{e_1} \cap \ldots \cap J_{e_n}, \qquad ext{s.t.} \qquad V(J_{e_i}) \subset U.$$

ldeal power, J^k computed by multiplying generators, e.g.

$$J^2 = \langle q_1, q_2, q_3 \rangle^2 = \langle q_1^2, q_2^2, q_3^2, q_1q_2, q_2q_3, q_1q_3 \rangle.$$

	Modern Techniques		Applications	Conclusions
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From Geometry to an Ansatz

Ansatz Construction Algorithm, Sketched

1. Construct branches of varieties where two \mathcal{D}_k vanish.

$$\mathcal{D}_i = \mathcal{D}_j = 0 \qquad \longrightarrow \qquad \mathcal{V} = \{U_1, U_2, \ldots\}.$$

2. Sample near surface to determine degree of divergence.

$$U: \quad \mathcal{D}_{i} \sim \mathcal{D}_{j} \sim \epsilon \qquad \Rightarrow \qquad \mathcal{C}_{k} \sim \frac{1}{\epsilon^{\kappa_{U}}}.$$

3. Ansatz is basis of intersection of associated ideals of vanishing polynomials. Constructed using Gröbner basis techniques.

$$\mathcal{N}_k \in \bigcap_{U \in \mathcal{V}} I(U)^{\langle \kappa_U \rangle}.$$

Studying a Six-Point Tree

$$\mathcal{A} = rac{\mathcal{N}}{\langle 12
angle \langle 23
angle \langle 34
angle [45] [56] [61] s_{345}}.$$



e.g.
$$[12] \sim [13] \sim [23] \sim \mathcal{O}(\epsilon) \qquad \Rightarrow \qquad \mathcal{A} \sim \mathcal{O}(\epsilon^2).$$

 $\begin{array}{l} \blacktriangleright \ \mathcal{N} \ \text{vanishes non-trivially on 28 varieties} \Rightarrow \text{ideal memberships:} \\ \mathcal{N} \ \in \ \left< [12], [13], [23] \right>^2 \cap \left< \langle 12 \rangle, \langle 34 \rangle \right> \cap \left< \langle 12 \rangle, [16] \right> \cap (25 \text{ more}). \end{array}$

Imposing mass dimension constraints gives one term Ansatz:

$$\mathcal{N} = \textbf{c}_0 \Big(\langle 12 \rangle [21] \langle 45 \rangle [54] \langle 4|2+3|1] \rangle + [16] \langle 6|1+2|3] \langle 34 \rangle \textbf{s}_{123} \Big), \quad \textbf{c}_0 \in \mathbb{Q}.$$

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Proof-of-Concept Remainders for $q\overline{q} \rightarrow \gamma\gamma\gamma$



(Analytic input from [Abreu, Page, Pascual, Sotnikov '20]).

• Codimension two surface study requires 317 \mathbb{Q}_p samples.

Fitting Ansatz now requires at most 566 \mathbb{F}_p samples.

Amplitude	$R^{(2,0)}_{-++}$	$R_{-++}^{(2,N_f)}$	$R^{(2,0)}_{+++}$	$R_{+++}^{(2,N_f)}$
Old Ansatz Dim	41301	2821	7905	1045
New Ansatz Dim	566	20	18	6

Introduction	Modern Techniques	Geometry of \mathcal{C}_k	Applications	Conclusions
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Summary

▶ Large progress in $2 \rightarrow 3$ scattering, including phenomenology!

Algebraic geometry provides insight into rational coefficients.

▶ For $q\overline{q} \rightarrow \gamma\gamma\gamma$ now only $\mathcal{O}(1000)$ numerical samples required.

From Ideals To Ansätze

 \blacktriangleright \mathcal{N} is a polynomial of dimension d and little group weight $\vec{\phi}$:

$$\mathcal{M}_{d,\phi} = \Big\{ \mathbf{a} \in \mathcal{R}_{\mathbf{n}} : [\mathbf{a}] = d, \{\mathbf{a}\}_{j} = \phi_{j} \Big\}.$$

• Construct independent spinor monomials $m_{\alpha,\beta}$ via Groebner.

▶ Numerator in each ideal. Intersect with $\mathcal{M}_{d,\vec{\phi}}$ (by Groebner).

$$\mathcal{M}_{d,\vec{\phi}} \cap \left[\bigcap_{U \in \mathcal{V}^{(2)}} I(U)^{\langle k_U \rangle}\right] = \bigcap_{U \in \mathcal{V}^{(2)}} \left[\mathcal{M}_{d,\vec{\phi}} \cap I(U)^{\langle k_U \rangle}\right]$$

Basis of intersection is Ansatz incorporating all constraints.

Massless 5-Point 2-Loop Singular Varieties (i)

▶ $\mathcal{V}^{(2)}$. Built from 10 parity/permutation generators:

Vi	1	2	3	4	5	6	7	8	9	10
# Perms	20	10	2	30	10	60	120	15	30	20

$$\begin{split} I(V_1) &= \langle \langle 12 \rangle, \langle 13 \rangle, \langle 23 \rangle, [45] \rangle_{R_5} ,\\ I(V_2) &= \langle \lambda_1^{\alpha} \rangle_{R_5} ,\\ I(V_3) &= \langle \langle 12 \rangle, \langle 23 \rangle, \langle 34 \rangle, \langle 45 \rangle, \langle 15 \rangle, \langle 13 \rangle, \langle 14 \rangle, \langle 24 \rangle, \langle 25 \rangle, \langle 35 \rangle \rangle_{R_5} ,\\ I(V_4) &= \langle \langle 12 \rangle, \langle 34 \rangle, \lambda_1^{\alpha} [15] + \lambda_2^{\alpha} [25] \rangle_{R_5} ,\\ I(V_5) &= \langle \langle 12 \rangle, [12] \rangle_{R_5} , \end{split}$$

Massless 5-Point 2-Loop Singular Varieties (ii)

$$I(V_6) = \langle \langle 12 \rangle, [13] \rangle_{R_5},$$

 $I(V_7) = \langle \langle 12 \rangle, \langle 3|1+4|3], \lambda_1^{lpha}[14][35] + \lambda_2^{lpha}[25][34]
angle_{R_5},$

$$\begin{split} I(V_8) &= \langle \langle 1|2+3|1], \langle 1|2+4|1], [1|2|4|1] - [1|2|3|1] + [1|3|4|1], \lambda \leftrightarrow \tilde{\lambda}, \\ &\langle 45 \rangle \langle 12 \rangle \langle 13 \rangle - \langle 15 \rangle \langle 23 \rangle \langle 14 \rangle, \lambda \leftrightarrow \tilde{\lambda} \rangle_{R_5} \,, \end{split}$$

$$egin{aligned} &I(V_9) = \langle \langle 1|2+3|1], \langle 2|1+3|2],\ & ilde{\lambda}_1^{\dotlpha} \langle 12
angle \langle 13
angle - ilde{\lambda}_2^{\dotlpha} \langle 12
angle \langle 23
angle - ilde{\lambda}_3^{\dotlpha} \langle 13
angle \langle 23
angle, \lambda \leftrightarrow ilde{\lambda}
angle_{R_5}, \end{aligned}$$

$$\begin{split} I(V_{10}) &= \langle \langle 1|2+3|1], \langle 2|1+4|2], \\ \lambda_1^{\alpha} [13] [14] [25] + \lambda_2^{\alpha} [12] [24] [35] + \lambda_3^{\alpha} [13] [24] [35], \lambda \leftrightarrow \tilde{\lambda} \rangle_{R_5} \,. \end{split}$$

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Irreducible Decomposition Examples

The once singular variety at four point is reducible

$$V\big(\langle \langle 12 \rangle \rangle_{R_4}\big) = V\big(\langle \langle 12 \rangle, [34] \rangle_{R_4}\big) \cup V\big(\langle \langle 12 \rangle, \langle 13 \rangle, \langle 14 \rangle, \langle 23 \rangle, \langle 24 \rangle, \langle 34 \rangle \rangle_{R_4}\big) .$$

▶ Many doubly singular varieties are reducible at five point, e.g. $V(\langle \langle 12 \rangle, \langle 34 \rangle \rangle_{R_5}) = V(\langle \langle 12 \rangle, \langle 34 \rangle, \lambda_1^{\alpha}[15] + \lambda_2^{\alpha}[25] \rangle_{R_5}) \cup V_{\text{angle}},$ $V(\langle \langle 12 \rangle, \langle 13 \rangle \rangle_{R_5}) = V(\langle \langle 12 \rangle, \langle 13 \rangle, \langle 23 \rangle, [45] \rangle_{R_5}) \cup V(\langle \lambda_1^{\alpha} \rangle_{R_5}) \cup V_{\text{angle}},$ where

$$V_{\mathsf{angle}} = V\big(\langle \langle 12 \rangle, \langle 23 \rangle, \langle 34 \rangle, \langle 45 \rangle, \langle 15 \rangle, \langle 13 \rangle, \langle 14 \rangle, \langle 24 \rangle, \langle 25 \rangle, \langle 35 \rangle \rangle_{\mathcal{R}_5} \big) \,.$$

Lorentz Invariance

Coefficients are Lorentz invariant functions of spinor brackets.

$$\mathcal{C}(\lambda, ilde{\lambda}) = \mathcal{C}(\langle
angle, []).$$

• Relevant ring is Lorentz invariant subring of S_n .

$$S_n = \mathbb{F}\Big[\langle 12 \rangle, \ldots, \langle (n-1)n \rangle, [12], \ldots [(n-1)n]\Big].$$

Variables are brackets, now have "Schouten identities".

$$\mathcal{J}_{\Lambda_n} = \left\langle \sum_{j=1}^n \langle ij \rangle [jk], \langle ij \rangle \langle kl \rangle - \langle ik \rangle \langle jl \rangle - \langle il \rangle \langle kj \rangle, \langle \rangle \leftrightarrow [] \right\rangle.$$

Physical spinor bracket functions also form a quotient ring.

$$\mathcal{R}_n = \mathcal{S}_n / \mathcal{J}_{\Lambda_n}.$$

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Bases of Spinor Space and Polynomial Reduction

▶ Numerators are Q−linear combinations of spinor monomials.

$$m_{\alpha} = \prod_{i} v_{i}^{\alpha_{i}}$$
 where $\vec{v} = \{\langle 12 \rangle, \langle 23 \rangle, \dots [12], [23], \dots \}.$

• Little group/mass dimension constraints pick finite set of α .

- Polynomial reduction writes p in terms of generators g_i . $p = \Delta_{\{g_1,...,g_k\}}(p) + \sum_{i=1}^k c_i g_i.$
- Polynomial in ideal if and only if Groebner remainder is 0.

$$\Delta_{\mathcal{G}(J)}\left(p
ight)=0\qquad \Leftrightarrow\qquad p\in J.$$

► Monomials irreducible by $\mathcal{G}(\mathcal{J}_{\Lambda_n})$ form basis. Related [Zhang '12] basis = { m_α such that $\Delta_{\mathcal{G}(\mathcal{J}_{\Lambda_n})}(m_\alpha) = m_\alpha$ }.

How To Perform Numerical Investigations?

• Need to find phase-space points $(\lambda^{\epsilon}, \tilde{\lambda}^{\epsilon})$ where \mathcal{D}_i are small.

$$\mathcal{D}_i(\lambda^\epsilon, \tilde{\lambda}^\epsilon) \sim \mathcal{D}_j(\lambda^\epsilon, \tilde{\lambda}^\epsilon) \sim \epsilon.$$

• Conflict with modern techniques: no small elements in \mathbb{F}_p .

$$|0|_{\mathbb{F}_p} = 0,$$
 and $a \neq 0 \Rightarrow |a|_{\mathbb{F}_p} = 1.$

Approaching with complex numbers plagued by instabilities.

Enter the p-adic numbers – a middle ground between \mathbb{F}_p and \mathbb{C} .

Introduction to the *p*-adic Numbers

The p-adic numbers roughly correspond to Laurent series in p.

$$x = \sum_{i=\nu}^{\infty} a_i p^i = a_{\nu} p^{\nu} + a_{\nu+1} p^{\nu+1} + \cdots, \qquad \begin{pmatrix} a_i \in [0, p-1], \\ a_{\nu} \neq 0. \end{pmatrix}.$$

▶ The *p*-adic numbers form a field. $x, y \in \mathbb{Q}_p \Rightarrow$

$$x+y\in \mathbb{Q}_p, \quad -x\in \mathbb{Q}_p, \quad x\times y\in \mathbb{Q}_p, \quad \frac{1}{x}\in \mathbb{Q}_p \text{ (if } x\neq 0\text{)}.$$

• The *p*-adic absolute value allows for small numbers $(p \sim \epsilon)$.

$$|x|_{
ho}=
ho^{-
u}, \quad \Rightarrow \quad |p|_{
ho}< |1|_{
ho}.$$

Computing with *p*-adic Numbers

For computing purposes* we truncate to finite order.

$$x = p^{\nu(x)} \Big(\underbrace{\tilde{x}}_{\text{mantissa}} + \mathcal{O}(p^k)\Big).$$

*Try [https://github.com/GDeLaurentis/pyadic] to investigate yourselves.

• Truncation reduces to finite field case for $\nu = 0, k = 1$.

• Arithmetic (+ - /*) is essentially performed modulo p^k , e.g.

$$x \times y = p^{\nu(x) + \nu(y)} \left(\tilde{x} \tilde{y} + \mathcal{O}(p^k) \right).$$

Mantissa inverse computed with extended euclidean algorithm.

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P-adic (Integer) Points Near an Irreducible Variety

▶ Want to find $(\lambda^{(\epsilon)}, \tilde{\lambda}^{(\epsilon)})$ "close" to $U = V(\langle q_1, \dots, q_m \rangle_{R_n})$:

$$q_i\left(\lambda^{(\epsilon)}, \tilde{\lambda}^{(\epsilon)}
ight) = pc_i + \mathcal{O}(p^k), \qquad \sum_{i=1}^n \lambda^{(\epsilon)}_{i\alpha} \tilde{\lambda}^{(\epsilon)}_{i\dot{lpha}} = 0 + \mathcal{O}(p^k).$$

First, find finite field $x \in U$ by intersecting with random plane.



• Arbitrarily extend \mathbb{F}_p point $(\lambda, \tilde{\lambda})$ to k digits. Trivially near U.

• To satisfy momentum conservation, perturb by $(p\delta, p\tilde{\delta})$.

$$(\lambda^{(\epsilon)}, \tilde{\lambda}^{(\epsilon)}) = (\lambda + p\delta, \tilde{\lambda} + p\tilde{\delta}).$$

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Examples of Symbolic Powers

► A function vanishing to fourth order at a point on the circle:

$$\langle x-1 \rangle_{\mathbb{F}[x,y]/\langle x^2+y^2-1 \rangle}^{\langle 4 \rangle} \sim$$

Often the symbolic power coincides with standard power, e.g.

$$\langle \langle 12 \rangle, [12] \rangle_{R_5}^{\langle 2 \rangle} = \langle \langle 12 \rangle, [12] \rangle_{R_5}^2 = \langle \langle 12 \rangle^2, \langle 12 \rangle [12], [12]^2 \rangle_{R_5}.$$

Symbolic/standard power may not coincide. E.g. in $\mathbb{F}[x, y, z]$ $\langle xy, xz, yz \rangle^{\langle 2 \rangle} = \langle x^2y^2, x^2z^2, y^2z^2, xyz \rangle \neq \langle xy, xz, yz \rangle^2$